# Calorimeter simulation, channel recovery, clustering

# Introductory talk



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Outline

► introduction

► G3-G4 calorimeter simulation

▷ highlights

dead channel recovery

▷ motivation

clustering algorithm

 $\triangleright$  introduction

► summary – future planning

# **Introducing GM**

### ► previously

: on the development and design of a tungsten/quartz fiber calorimeter for the ALICE experiment at the CERN-LHC

#### ▶ since 030801

: CALICE Collaboration

#### ▶ in this talk

- : highlight some first results
- : outline fields of interest and work
- : feedback is welcome

# **G3-G4 calorimeter simulation**

- ▶.
- : work with the latest versions of BRAHMS(v3.01) and MOKKA(v2.0) porting calorimeter geometry from G4
- : analysis code in ROOT and C++

### : reproduce older results, cross-check

: develop Event class - tree, data in ascii grow fast, possibly develop "root2hbook" or "root2txt" backward utility



# **G3-G4** comparison



E deposited per HCAL cell







E deposited per ECAL cell

Number of cells hit vs layer



# **G3-G4** comparison

counts

160

140

120

100

80



#### E deposited per ECAL cell









#### E deposited per HCAL cell

1000 e<sup>-</sup> 5 GeV

G4 (RangeCut 5 µm)

G3 (EnergyCut 10 keV)

E deposited (MeV)

response vs G3 control parameter (energy cut)



# response vs G4 control parameter (range cut)



resolution vs G3 control parameter (energy cut)



# resolution vs G4 control parameter (range cut)



# width vs G3 control parameter (energy cut)



# width vs G4 control parameter (range cut)



# **CPU time vs G3 control parameter (energy cut)**



# **CPU time vs G4 control parameter (range cut)**





# **Dead channel recovery**

#### ► fact

: production cost considerations may favour dead channel fraction around 5%

#### ► questions

- : how this affects performance
- : can we compensate recover dead channels

#### ▶ study - solution

- : study detector performance vs dead channel fraction
- : develop dead channel recovery scheme

besides conventional methods, potential application of neural network based techniques

: evaluate recovery efficiency and performance improvement



# **Illustration of the problem**



Figure 1: original picture, 10000 channels



Figure 2: picture with 5% dead channels

# **Illustration of the solution**



Figure 3: "recovery scheme" step 1



Figure 4: "recovery scheme" step 2

# **Illustration of the solution**



Figure 5: "recovery scheme" step 3



Figure 6: recovered picture



•  $I_i^k$  : total input to neuron i of layer k

$$I_i^k = \sum_{j=1}^{N^{k-1}} w_{ji}^{k-1} O_j^{k-1} \quad , \quad k > 1$$
<sup>(1)</sup>

•  $O_i^k$  : neuron output,  $w_{ji}^{k-1}$  : connection weight with neuron j of layer (k-1)

$$O_i^k = f(I_i^k + w_{i0}^k) , \ k > 1$$
 (2)

- $f(x) \equiv$  sigmoid activation function, e.g.  $f(x) = \frac{1}{1 + e^{-x/T}}$
- number of free parameters to be optimized

$$N_{weights} = \sum_{k=2}^{H+2} (N^k + N^k N^{k-1})$$
(3)





Figure 7: typical signal and background distribution as a function of NNoutput. The hatched area contains the  $N_{signalNN}$  signal events that are above the selection cut NNoutput<sub>cut</sub>. The colored area contains the contaminating  $N_{signallikeNN}$  background event s.

signal efficiency: 
$$\epsilon_s = \frac{N_{signalNN}}{N_{signal}}$$
 (4)

contamination: 
$$\epsilon_b = \frac{N_{signallikeNN}}{N_{background}}$$
 (5)

signal enhancement = 
$$\frac{\epsilon_s}{\epsilon_b}$$
 (6)

$$(S/B)_{NN} = \frac{N_{signalNN}}{N_{signallikeNN}} = \frac{\epsilon_s}{\epsilon_b} \cdot \frac{N_{signal}}{N_{background}} = \frac{\epsilon_s}{\epsilon_b} \cdot S/B$$
(7)



# **Clustering algorithm**

#### minimal spanning tree

: a tree which contains all nodes with no circuits and of which the sum of weights of its edges is minimum

#### ► properties

- : unique for the given set of nodes and the chosen metric
- : deterministic, no dependence on random choices of nodes
- : invariant under similarity transformations that preserve the monotony of the metric

# MST and clustering

- : theorem 1: any MST contains at least one edge from each link-set between P and Q partitions
- : theorem 2: all MST edges are links of some partition of graph
- : theorem 3: if S denotes the nodes of graph and C is a non-empty subset of S with the property that ρ(P,Q) < ρ(C,S-C) for all partitions P, Q of C, then the restriction of any MST to the nodes of C forms a connected subtree of the MST
- : *theorem 4:* if T is an MST for graph G and X, Y are two nodes of G then the unique path in T from X to Y is a minimax path from X to Y

# References

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# **Summary – Future planning**

# ► G3-G4 calorimeter simulation

- : high priority
- : make cut / particle / energy scan with G3 and G4
- : studies to concentrate on prototype geometry
- : identify regions where testbeam should focus to give answers

### ► channel recovery

- : study detector performance vs dead channel fraction
- : develop dead channel recovery scheme
- : evaluate recovery efficiency and performance improvement

# clustering algorithm

: develop clustering algorithm based on MST approach

